

Validated simulation of DAEs Julien Alexandre dit Sandretto





Department U2IS ENSTA Paris SSC310-2020 Contents



Guaranteed simulation of differential equations

Differential Algebraic Equations

Approach to simulate DAE

Examples

Recall of Ordinary differential equations



Given by

$$y'=f(y,t)$$

Initial Value Problems

$$y'=f(y,t), \quad y(0)=y_0$$

Numerical simulation of IVPs till a time t_n Compute $y_i \approx y(t_i)$ with $t_i \in \{0, t_1, \dots, t_n\}$

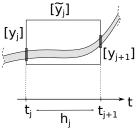
Validated simulation of IVPs

Produces a list of boxes $[y_i]$ and $[\tilde{y}_i]$ such that

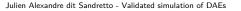
- $y(t_i) \in [y_i]$ with $t_i \in \{0, t_1, ..., t_n\}$
- $y(t) \in [\tilde{y}_i]$ for all $t \in [t_i, t_{i+1}]$

Method of Lohner

- 1. Find $[\tilde{y}_i]$ with Picard-Lindelof operator
- Compute [y_i] with a validated integration scheme : Taylor (Vnode-LP) or Runge-Kutta (DynIbex)









Differential Algebraic Equations

General form: implicit

 $F(t, y, y', ...) = 0, t_0 \le t \le t_{end}$ $y' = \mathsf{DAE} \ 1^{st} \text{ order}, y'' = \mathsf{DAE} \ 2^{nd}, \text{ etc.}$ (all DAEs can be rewritten in DAE of 1^{st} order)



Hessenberg form: Semi-explicit (index: distance to ODE)

index 1:
$$\begin{cases} y' = f(t, x, y) \\ 0 = g(t, x, y) \end{cases}$$

$$\left(\text{index } 2: \begin{cases} y' = f(t, x, y) \\ 0 = g(t, x) \end{cases} \right)$$

 \Rightarrow Focus on Hessenberg index-1: Simulink, Modelica-like, etc.

Different from ODE + constraint

$$y' = f(t, y)$$

 $0 = g(y, y')$, $t_0 \le t \le t_{end}$
 \Rightarrow Direct with contractor approximately $f(t, y)$

Differential Algebraic Equations

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Hessenberg form: Semi-explicit (index: distance to ODE)

$$(index \ 1: \left\{ \begin{array}{c} y' = f(t, x, y) \\ 0 = g(t, x, y) \end{array} \right) \qquad (index \ 2: \left\{ \begin{array}{c} y' = f(t, x, y) \\ 0 = g(t, x) \end{array} \right)$$

Some of dependent variables occur without their derivatives ! Different from ODE + constraint

$$\left\{ \begin{array}{l} y' = f(t,y) \\ 0 = g(y,y') \end{array} \right., \ t_0 \leq t \leq t_{end} \\ \Rightarrow \ \mathsf{Direct} \ \mathsf{with} \ \mathsf{contractor} \ \mathsf{approach} \end{array} \right.$$



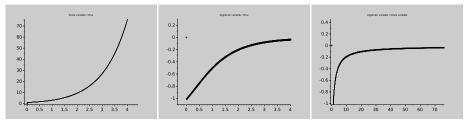
A basic example



System in Hessenberg index-1 form

$$\begin{cases} y' = y + x + 1\\ (y+1) * x + 2 = 0 \end{cases} \quad y(0) = 1.0 \text{ and } x(0) = 0.0$$

Simulation \Rightarrow stiffness (in general)



Simulation of a DAE

As ODE: a list of boxes $[y_i]$ and $[\tilde{y}_i]$ such that

- $y(t_i) \in [y_i]$ with $t_i \in \{0, t_1, \ldots, t_n\}$
- $y(t) \in [\tilde{y}_i]$ for all $t \in [t_i, t_{i+1}]$

But in addition: a list of boxes $[x_i]$ and $[\tilde{x}_i]$ such that

- $x(t_i) \in [x_i]$ with $t_i \in \{0, t_1, ..., t_n\}$
- $x(t) \in [\tilde{x}_i]$ for all $t \in [t_i, t_{i+1}]$

Both validate

- ▶ $y'(t_i) \in f(t_i, [x_i], [y_i])$
- $\blacktriangleright \exists x \in [x_i], \exists y \in [y_i] : g(t_i, x, y) = 0$
- ► $y'(t) \in f(t, [\tilde{x}_i], [\tilde{y}_i]), \forall t \in [t_i, t_{i+1}]$
- $\forall t \in [t_i, t_{i+1}], \exists x \in [\tilde{x}_i], \exists y \in [\tilde{y}_i] : g(t, x, y) = 0$

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- $\flat \ \forall t \in [t_i, t_{i+1}], \exists x \in [\tilde{x}_i], \exists y \in [\tilde{y}_i] : g(t, x, y) = 0$





Step 1- A priori enclosure of state and algebraic variables

How find the enclosure $[\tilde{x}]$ on integration step ?

Assume that $\frac{\partial g}{\partial x}$ is locally reversal

we are able to find the unique $x = \psi(y)$ (implicit function theorem), and then:

$$y'=f(\psi(y),y)$$

and finally we could apply Picard-Lindelof to prove **existence and uniqueness**, but...



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 ψ is unknown !



Step 1- A priori enclosure of state and algebraic variables

Solution

If we are able to find $[\tilde{x}]$ such that for each $y \in [\tilde{y}], \exists ! x \in [\tilde{x}] : g(x, y) = 0$, then $\exists ! h$ on the neighborhood of $[\tilde{x}]$, and the solution of DAE $\exists !$ in $[\tilde{y}]$ (Picard with $[\tilde{x}]$ as a parameter)

A novel operator Picard-Krawczyk \mathcal{PK} : If $\begin{pmatrix} \mathcal{P}([\tilde{y}], [\tilde{x}]) \\ \mathcal{K}([\tilde{y}], [\tilde{x}]) \end{pmatrix} \subset Int \begin{pmatrix} [\tilde{y}] \\ [\tilde{x}] \end{pmatrix}$ then $\exists !$ solution of DAE

- \mathcal{P} a Picard-Lindelof for $y' \in f([\tilde{x}], y)$
- \mathcal{K} a parametrized preconditioned Krawczyk operator for $g(x,y) = 0, \forall y \in [\tilde{y}]$

Parametric Krawczyk



Parametric preconditioned Krawczyk operator

$$\mathcal{K}([\tilde{y}], [\tilde{x}]) = m([\tilde{x}]) - Cg(m([\tilde{x}]), m([\tilde{y}])) - (C\frac{\partial g}{\partial x}([\tilde{x}], [\tilde{y}]) - I)([\tilde{x}] - m([\tilde{x}])) - C\frac{\partial g}{\partial y}(m([\tilde{x}]), [\tilde{y}])([\tilde{y}] - m([\tilde{y}]))$$
(1)

Parametric Krawczyk



Interval Newton operator $\mathcal{N}([x])$:

repeat

$$\begin{split} & [A] = J([x]) \\ & [b] = F(\mathrm{m}([x])) \\ & \text{Solve } [A]s = [b] \text{ with a linear system solver method (Gauss elimination for example)} \\ & [x] = [x] \cap s + \mathrm{m}([x]) \\ & \text{until Fixed point} \end{split}$$

If $\mathcal{N}([x]) \subset Int([x])$, then F has a unique solution and this solution is in $\mathcal{N}([x])$

Parametric preconditioned Krawczyk

A better version of Newton, with parameter and preconditioning

Frobenius theorem



Let X and Y be Banach spaces, and $A \subset X$, $B \subset Y$ a pair of open sets. Let $F : A \times B \to L(X, Y)$

be a continuously differentiable function of the Cartesian product (which inherits a differentiable structure from its inclusion into $X \times Y$) into the space L(X,Y) of continuous linear transformations of X into Y. A differentiable mapping $u : A \rightarrow B$ is a solution of the differential equation v' = F(x, v) (1)

If u'(x) = F(x, u(x)) for all $x \in A$. The equation (1) is completely integrable if for each $(x_0, y_0) \in A \times B$, there is a neighborhood U of x0 such that (1) has a unique solution u(x) defined on U such that $u(x_0)=y_0$. The conditions of the Frobenius theorem depend on whether the underlying field is R or C. If it is R, then assume F is continuously differentiable. If it is C, then assume F is twice continuously differentiable. Then (1) is completely integrable at each point of $A \times B$ if and only if

 $\begin{array}{l} D_1F(x,y) \cdot (s_1,s_2) + D_2F(x,y) \cdot (F(x,y) \cdot s_1,s_2) = D_1F(x,y) \cdot (s_2,s_1) + D_2F(x,y) \cdot (F(x,y) \cdot s_2,s_1) \text{ for all } s_1,s_2 \in X. \text{ Here D1} (\text{resp. D2}) \text{ denotes the partial derivative with respect to the first (resp. second) variable; the dot product denotes the action of the linear operator <math>F(x,y) \in L(X,Y)$, as well as the actions of the operators $D_1F(x,y) \in L(X,L(X,Y))$ and $D_2F(x,y) \in L(Y,L(X,Y))$.

Dieudonné, J (1969). Foundations of modern analysis. Academic Press. Chapter 10.9.



Step 2- Contraction of state and algebraic variables (at t + h)

Two contractors in a fixpoint:

- Contraction of $[y_{i+1}]$ (init $[\tilde{y}_i]$)
 - $[\tilde{x}_i]$ as a parameter of function f(t, x, y)
 - \Rightarrow ODE (stiff + interval parameter)
 - \Rightarrow Radau IIA order 3 (fully Implicit Runge-Kutta, A-stable, efficiency for stiff and interval parameters)
- Contraction of $[x_{i+1}]$ (init $[\tilde{x}_i]$)
 - $[y_{i+1}]$ as a parameter of function g(x, y)

 $\Rightarrow \mathsf{Constraint} \ \mathsf{solving}$

 $\Rightarrow \mathsf{Krawczyk} + \mathsf{forward}/\mathsf{backward}$

(+ any other constraints, from physical context or Pantelides algorithm)

Recall on Radau methods



$$y_{n+1} = y_n + h \sum_{i=1}^{s} b_i k_i, \quad k_i = f\left(t_0 + c_i h, y_0 + h \sum_{j=1}^{s} a_{ij} k_j\right)$$

Butcher tableau Radau IIA order 3

Butcher tableau Radau IIA order 5

Approach to simulate DAE

Based on Lohner two-step approach



How to control the stepsize of integration scheme ? Classical method: Constrained by the Picard success and an evaluation of the truncature error lower than threshold

No specific control w.r.t. the algebraic variable If x leads to a large evaluation of truncature error: too late !

Solution: force diameter of x grows slower than yEmpirical approach: to improve ! Approach to simulate DAE

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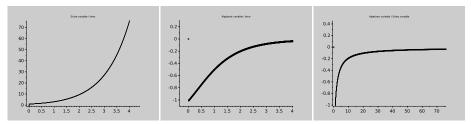
Solution: force diameter of x grows slower than yEmpirical approach: to improve ! A basic example



System in Hessenberg index-1 form

$$\begin{cases} y' = y + x + 1\\ (y+1) * x + 2 = 0\\ (\text{consistency: } x(0) = -1) \end{cases} y(0) = 1.0 \text{ and } x(0) \in [-2.0, 2.0]$$

Simulation till t=4s (30 seconds of computation)



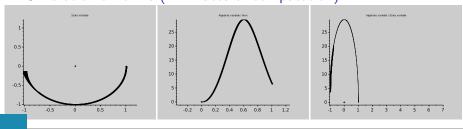
Examples

The classical example: Pendulum



$$\begin{cases} p' = u \\ q' = v \\ mu' = -p\lambda \\ mv' = -q\lambda - g \end{cases}$$
$$m(u^{2} + v^{2}) - gq - l^{2}\lambda = 0$$

 $(p, q, u, v)_0 = (1, 0, 0, 0)$ et $\lambda_0 \in [-0.1, 0.1]$ (consistency: $\lambda = 0$) Simulation till t=1s (2 minutes of computation)



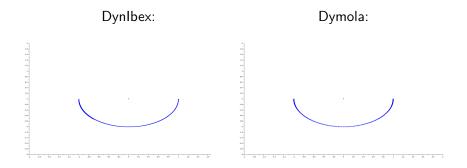
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Examples

Pendulum with Dymola





Pantelides on pendulum



Pantelides

Algorithm for order reduction, formal differentiation and manipulation of equations

On pendulum problem

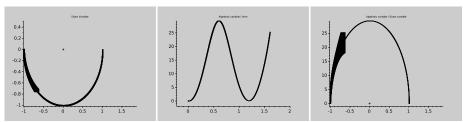
$$\begin{cases} p^2 + q^2 - l^2 = 0\\ p * u + q * v = 0\\ m * (u^2 + v^2) - g * q^2 - l^2 * p = 0 \end{cases}$$

 \Rightarrow Constraints valid all the time !

Examples

Pendulum to 1.6*s*, $tol = 10^{-18}$

28 minutes...



With csp: 27 minutes...

